GlobalFS: A Strongly Consistent Multi-Site Filesystem

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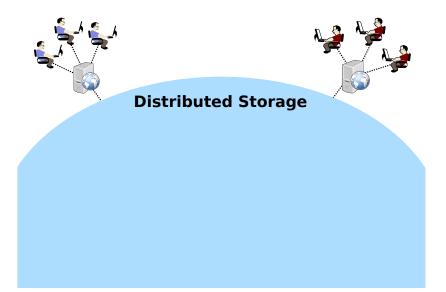


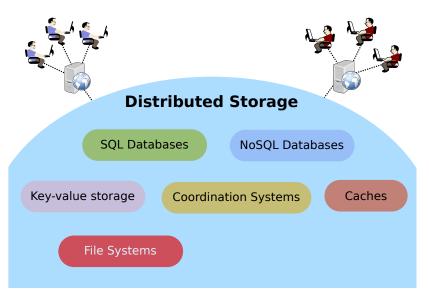


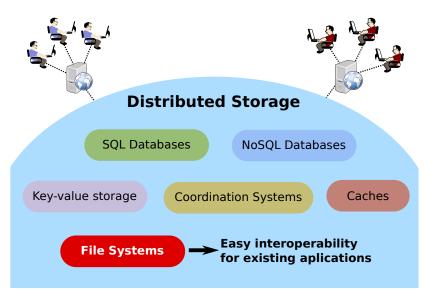












Global infrastructure



Amazon's AWS global infrastructure

CAP theorem

Weak Consistency

Lower latency

Higher availability

Possibly incorrect/unexpected results

Strong Consistency

Clear semantics and guarantees

Easier to reason about

Block instead of providing incorrect results

What is GlobalFS?

Geographically distributed filesystem

Familiar interface (POSIX)

Strong consistency

Fault-tolerance through replication

Flexible performance through locality

Overall design

Separate data and metadata

Partial replication

Metadata protocol exploiting atomic multicast

Causal reads

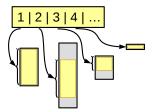
Separate data and metadata

Immutable data Variable sized blobs

Metadata

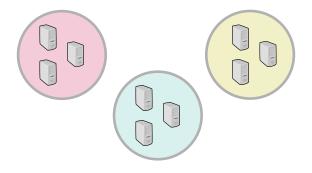
Controls file contents, properties and filesystem structure

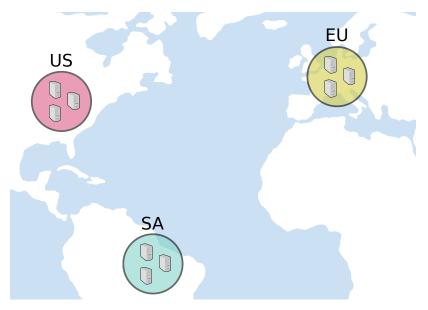
Metadata refers to data blobs

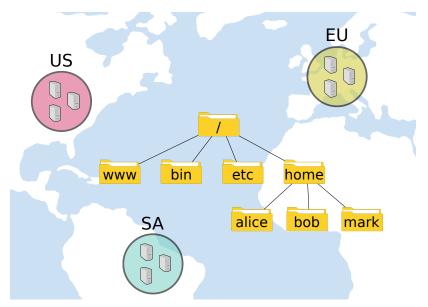


Immutable data is simple to replicate consistently

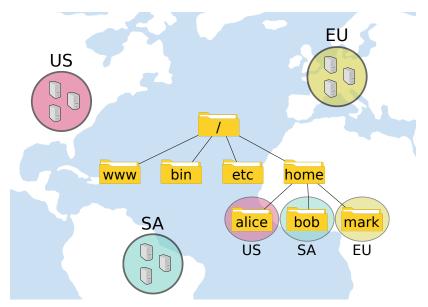
Metadata is partitioned between replica groups (i.e., partitions)

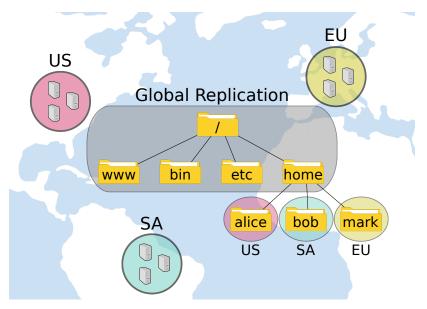




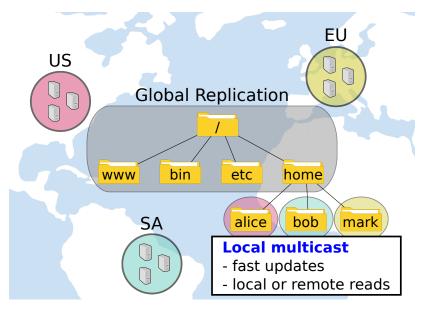


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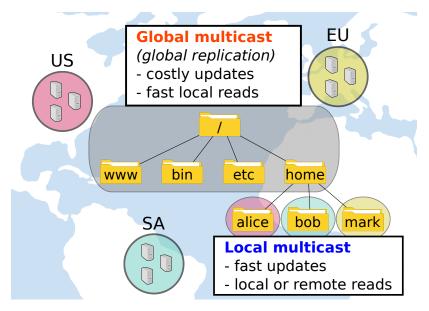




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Partial ordering

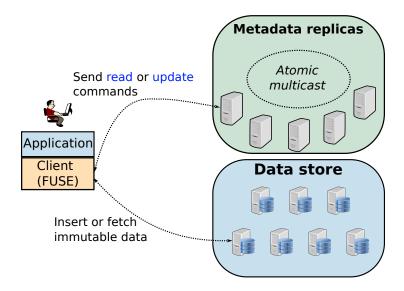
GlobalFS exploits atomic multicast

Atomic delivery to groups of processes

Partial ordering: messages for different groups don't have to be ordered betweem themselves

Partial ordering is critical for scalability

Architecture



Consistent update operations

Step 1 Write data blobs to data store Step 2 Issue a metadata update

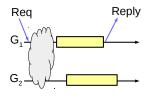
Consistent update operations

Step 1 Write data blobs to data store Step 2 Issue a metadata update

 $\begin{array}{c} \text{Req} \\ G_1 \\ \hline \\ G_2 \\ \hline \\ write to file in G_1 \end{array}$

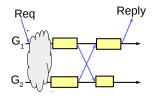
Single-partition

Uncoordinated multi-partition

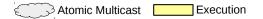


write to file in $\{G_1, G_2\}$





move file from G₁ to G₂



Causally related updates are seen in the same order *e.g., operations done by the same client*

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Client A

Creates an image cat.jpg

Modifies a page pets.html to include the image cat.jpg

Causally related updates are seen in the same order *e.g., operations done by the same client*

Client A

Client B

Creates an image cat.jpg

Modifies a page pets.html to include the image cat.jpg

Opens the pets.html page and finds a broken image reference

Where is the cat?

Step 1 Contact a metadata replica for a list of blob idsStep 2 Get the data from the data store

Approach inspired by vector clocks

Vector is composed of one counter per replica group

Evaluation

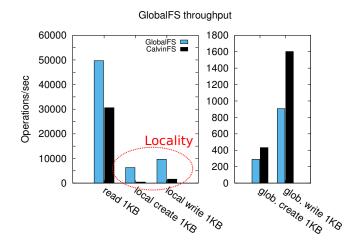
Complete prototype in Java https://github.com/pacheco/GlobalFS

Filesystem in Userspace (FUSE)

URingPaxos for atomic multicast

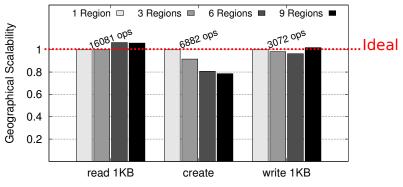
Global deployment using Amazon EC2

Maximum throughput by operation



3 region deployment US west, US east and Europe

Geographical scalability



Normalized throughput per region as more regions are added

9 regions uses all EC2 regions available at the time

GlobalFS: Summary

Strong consistency at global scale

Simple and familiar API (POSIX)

Flexible performance through partial replication and locality

Cheap causal read operations

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Thank you! Leandro Pacheco pachecol@usi.ch